

# Relational coloring of varieties containing a cube term

Matthew Moore

Vanderbilt University

May 21, 2016

# Coloring varieties containing a cube term

- 1 Coloring
- 2 Clone  $\mathcal{L}$ -homomorphisms
- 3 Linear interpretability
- 4 Further directions

# Coloring varieties containing a cube term

- 1 Coloring
- 2 Clone  $\mathcal{L}$ -homomorphisms
- 3 Linear interpretability
- 4 Further directions

# Cube terms

An  **$n$ -dimensional cube term** for an algebra  $\mathbb{A}$  is a term  $t(\dots)$  such that  $\mathbb{A} \models t(u_1, \dots, u_m) \approx \bar{x}$  for some  $u_i \in \{x, y\}^n \setminus \{x\}^n$ .

## Example

$$\mathbb{A} \models t \begin{pmatrix} y & x & x & y & y & x & y \\ x & y & x & y & x & y & y \\ x & x & y & x & y & y & y \end{pmatrix} \approx \begin{pmatrix} x \\ x \\ x \end{pmatrix}$$

**Finite idempotent**  $\mathbb{A}$  has a cube term...

$\Leftrightarrow \mathbb{A}$  has few subpowers:  $\log_2 |\mathbf{S}(\mathbb{A}^n)| \in \mathcal{O}(n^k)$ .

$\Leftrightarrow \mathbb{A}$  is congruence modular and finitely related.

$\Leftrightarrow \mathbb{A}$  has no cube term blockers:  $\mathbb{D} < \mathbb{C} \leq \mathbb{A}$  with  $C^n \setminus (C \setminus D)^n \leq \mathbb{A}^n$ .

$\Rightarrow \text{CSP}(\mathbb{A})$  is tractable.

# Coloring

Let  $\mathcal{A}$  be a clone and  $\mathbb{B} = (B; (R_j)_{j \in J})$  be a relational structure.

Let  $F^{\mathcal{A}}(B)$  be the free algebra in  $\mathcal{V}(\langle \mathcal{A}; \mathcal{A} \rangle)$  with generators  $B$ .

For  $R_i \in (R_j)_{j \in J}$  let

$$R_i^{\mathcal{A}} = \text{closure of } \{(b_1, \dots, b_n) \in F^{\mathcal{A}}(B)^n \mid (b_1, \dots, b_n) \in R_i\} \text{ under } \mathcal{A}$$

Let  $\mathbb{F}^{\mathcal{A}}(\mathbb{B}) = (F^{\mathcal{A}}(B); (R_j^{\mathcal{A}})_{j \in J})$ .

A **coloring** of  $\mathcal{A}$  by  $\mathbb{B}$  is a relational homomorphism

$$c : \mathbb{F}^{\mathcal{A}}(\mathbb{B}) \rightarrow \mathbb{B} \quad \text{such that} \quad c(b) = b.$$

We say that  $\mathcal{A}$  is  **$\mathbb{B}$ -colorable**. We say that  $\mathcal{V}$  is  **$\mathbb{B}$ -colorable** if  $\text{Clo}(\mathcal{V})$  is.

# What is it good for?

## Theorem (Sequeira)

$\mathcal{V}$  is congruence  $k$ -permutable for some  $k$  iff  $\mathcal{V}$  is not  $(\{0, 1\}; \leq)$ -colorable.

## Theorem (Sequeira)

Let  $D = \{1, 2, 3, 4\}$ ,  $\alpha = 12|34$ ,  $\beta = 13|24$ ,  $\gamma = 12|3|4$ , and  $\mathbb{D} = (D; \alpha, \beta, \gamma)$ .  $\mathcal{V}$  is congruence modular iff  $\mathcal{V}$  is not  $\mathbb{D}$ -colorable.

NB: Sequeira calls this “compatibility with projections”.

## Theorem

Let  $\emptyset \neq C \subsetneq B \subseteq A$  be sets,  $R_n = B^n \setminus (B \setminus C)^n$ , and  $\mathbb{A} = (A; (R_n)_{n \in \omega})$ .  
Idempotent  $\mathcal{V}$  has a cube term iff  $\mathcal{V}$  is not  $\mathbb{A}$ -colorable.

## Theorem

Let  $R_n = (\omega^2)^n \setminus (\Delta_2)^n$  and  $\mathbb{W} = (\omega; (R_n)_{n \in \omega})$ .  
 $\mathcal{V}$  has a weak Taylor term iff  $\mathcal{V}$  is not  $\mathbb{W}$ -colorable.

## Theorem

Let  $\emptyset \neq C \subsetneq B \subseteq A$  be sets,  $R_n = B^n \setminus (B \setminus C)^n$ , and  $\mathbb{A} = (A; (R_n)_{n \in \omega})$ .  
Idempotent  $\mathcal{V}$  has a cube term iff  $\mathcal{V}$  is not  $\mathbb{A}$ -colorable.

## Proof.

( $\Rightarrow$ ): Suppose  $\mathcal{V}$  has an  $n$ -dimensional cube term and is  $\mathbb{A}$ -colorable.

There is some identity  $\mathcal{V} \models t(u_1, \dots, u_n) \approx \bar{x}$  for some  $u_i \in \{x, y\}^n \setminus \{x\}^n$ .

Let  $c \in C$  and  $b \in B \setminus C$ .

Substitute  $b$  for  $x$  and  $c$  for  $y$  to get  $t(u_1, \dots, u_n) = \bar{b}$  and  $u_i \in R_n^{\mathcal{V}}$ .

Thus  $\bar{b} \in R_n^{\mathcal{V}}$ . Then  $c(\bar{b}) = \bar{b} \in R_n \subseteq A^n$ , a contradiction.



## Proof (cont.).

( $\Leftarrow$ ): Suppose  $\mathcal{V}$  does not have a cube term.

Write  $g \prec_t h$  if we have  $u_i \in \{g, h\}^n \setminus \{g\}^n$  such that  $t(u_1, \dots, u_n) = \bar{g}$ .

Fix some  $c_0 \in C$  and define  $c : \mathbb{F}^{\text{Clo}(\mathcal{V})}(A) \rightarrow \mathbb{A}$  by

$$c(f) = \begin{cases} a & \text{if } f = a \in A \subseteq F^{\text{Clo}(\mathcal{V})}(A), \\ b & \text{else if } \exists b \in B, t(\dots) \text{ such that } b \prec_t f, \\ c_0 & \text{else.} \end{cases}$$

Clearly  $c(a) = a$  for  $a \in A$ . Does  $c$  preserve relations?

## Proof (cont.).

$$c(f) = \begin{cases} a & \text{if } f = a \in A \subseteq F^{\text{Clo}(\mathcal{V})}(A), \\ b & \text{else if } \exists b \in B, t(\dots) \text{ such that } b \prec_t f, \\ c_0 & \text{else.} \end{cases}$$

Suppose that  $c$  fails to preserve  $R_n^{\text{Clo}(\mathcal{V})}$ . Then there exists

- $(f_1, \dots, f_n) \in R_n^{\text{Clo}(\mathcal{V})}$ ,
- a term  $s(\dots)$ ,
- and  $u_1, \dots, u_k \in B^n \setminus (B \setminus C)^n \subseteq F^{\text{Clo}(\mathcal{V})}(A)$ ,

such that

- $s(u_1, \dots, u_k) = (f_1, \dots, f_n)$
- and  $c(f_i) = b_i \in (B \setminus C)$ .

Thus  $b_i \prec_{t_i} f_i$ . Let  $t = t_1 * t_2 * \dots * t_n * s$ .

Then  $\exists v_1, \dots, v_m \in B^n \setminus (B \setminus C)^n$  such that  $t(v_1, \dots, v_m) = (b_1, \dots, b_n)$ .

This is a free algebra, so substitute  $x$  for  $B \setminus C$  and  $y$  for  $C$  to obtain a cube identity for  $t(\dots)$ , a contradiction. □

# A corollary

## Theorem

Let  $\emptyset \neq C \subsetneq B \subseteq A$  be sets,  $R_n = B^n \setminus (B \setminus C)^n$ , and  $\mathbb{A} = (A; (R_n)_{n \in \omega})$ .  
Idempotent  $\mathcal{V}$  has a cube term iff  $\mathcal{V}$  is not  $\mathbb{A}$ -colorable.

Let  $A = B = \{0, 1\}$  and  $C = \{0\}$ .

Then the polymorphism clone of  $\mathbb{A}$  is generated by  $\hat{\mathcal{D}} = \langle \{0, 1\}; \nearrow \rangle$ :  
 $\text{Pol}(\mathbb{A}) = \text{Clo}(\hat{\mathcal{D}})$ .

## Corollary

Idempotent  $\mathcal{V}$  has a cube term iff  $\mathcal{V}$  is not  $\text{Rel}(\hat{\mathcal{D}})$ -colorable.

# Coloring varieties containing a cube term

- 1 Coloring
- 2 Clone  $\mathcal{L}$ -homomorphisms
- 3 Linear interpretability
- 4 Further directions

# Clone homomorphisms

## Definition

A homomorphism between clones  $\mathcal{A}$  and  $\mathcal{B}$  is an arity-preserving mapping  $\varphi : \mathcal{A} \rightarrow \mathcal{B}$  such that

$$\varphi(\pi_k) = \pi_k \quad \text{and} \quad \varphi(f(g_1, \dots, g_n)) = \varphi(f)(\varphi(g_1), \dots, \varphi(g_n)).$$

An  $\mathcal{L}$ -homomorphism between clones  $\mathcal{A}$  and  $\mathcal{B}$  is an arity preserving mapping  $\varphi : \mathcal{A} \rightarrow \mathcal{B}$  such that

$$\varphi(\pi_k) = \pi_k \quad \text{and} \quad \varphi(f(\pi_{i_1}, \dots, \pi_{i_n})) = \varphi(f)(\pi_{i_1}, \dots, \pi_{i_n}).$$

Let  $\mathcal{A}$  and  $\mathcal{B}$  be clones, and let  $\varphi : \mathcal{A} \rightarrow \mathcal{B}$  be an arity preserving map.

Let  $\mathbb{A} = (A; \mathcal{A})$ , and  $\mathbb{B} = (B; \mathcal{B})$ .

$\varphi$  is a clone homomorphism iff  $f \approx g$  in  $\mathbb{A}$  implies  $\varphi(f) \approx \varphi(g)$  in  $\mathbb{B}$ .

$\varphi$  is a clone  $\mathcal{L}$ -homomorphism iff every identity  $f \approx g$  in  $\mathbb{A}$  not involving composition implies  $\varphi(f) \approx \varphi(g)$  in  $\mathbb{B}$ .

## Example

Consider  $\mathfrak{B} = \langle \{0, 1\}; \rightarrow \rangle$  and  $\hat{\mathfrak{B}} = \langle \{0, 1\}; \nrightarrow \rangle$ .

No clone homomorphisms between  $\text{Clo}(\hat{\mathfrak{B}})$  and  $\text{Clo}(\mathfrak{B})$  since  $\rightarrow$  cannot be defined as a term in  $\nrightarrow$  and  $\nrightarrow$  cannot be defined as a term in  $\rightarrow$ .

This is also witnessed by the identities

$$\mathfrak{B} \models y \rightarrow (x \rightarrow x) \approx (x \rightarrow x), \quad \hat{\mathfrak{B}} \models (x \nrightarrow x) \nrightarrow y \approx (x \nrightarrow x).$$

Given a function  $f : \{0, 1\}^n \rightarrow \{0, 1\}$  define

$\delta(f(x_1, \dots, x_n)) = \neg f(\neg x_1, \dots, \neg x_n)$ , where  $\neg$  is boolean negation.

$\delta : \text{Clo}(\mathfrak{B}) \rightarrow \text{Clo}(\hat{\mathfrak{B}})$  is a clone  $\mathcal{L}$ -homomorphism, but not a clone homomorphism.

Let  $s(x, y, z) = x \rightarrow (y \rightarrow z)$  and  $s'(x, y, z) = (x \nrightarrow y) \nrightarrow z$ .

$$\delta(s) = s' \text{ and } \delta(s') = s.$$

$\delta$  translates the identities witnessing the non-existence of a clone homomorphism.

## Theorem (Barto, Oprsal, Pinsker)

*Let  $\mathcal{A}$  be a clone and  $\mathbb{B}$  a relational structure with polymorphism clone  $\mathcal{B}$ .  $\mathcal{A}$  is  $\text{Rel}(\mathbb{B})$ -colorable iff there is a clone  $\mathcal{L}$ -homomorphism  $\mathcal{A} \rightarrow \mathcal{B}$ .*

## Corollary

*Let  $\mathcal{V}$  be an idempotent variety. The following are equivalent.*

- $\mathcal{V}$  does not have a cube term,*
- there is a clone  $\mathcal{L}$ -homomorphism  $\text{Clo}(\mathcal{V}) \rightarrow \text{Clo}(\hat{\mathcal{2}})$ .*

# Coloring varieties containing a cube term

- 1 Coloring
- 2 Clone  $\mathcal{L}$ -homomorphisms
- 3 Linear interpretability
- 4 Further directions



# Interpretability

Let  $\mathcal{V}$  and  $\mathcal{W}$  be varieties and write  $\mathcal{V} \leq \mathcal{W}$  if there is a clone homomorphism  $\text{Clo}(\mathcal{V}) \rightarrow \text{Clo}(\mathcal{W})$ . Say that  $\mathcal{V}$  is **interpretable into**  $\mathcal{W}$ .

Define an equivalence relation by  $\mathcal{V} \equiv \mathcal{W}$  if  $\mathcal{V} \leq \mathcal{W} \leq \mathcal{V}$ .

Modulo this equivalence relation, under the  $\leq$  order the class of all varieties is a lattice. This is the **lattice of interpretability types**.

Write  $\mathcal{V} \leq_{\mathcal{L}} \mathcal{W}$  if there is a clone  $\mathcal{L}$ -homomorphism  $\text{Clo}(\mathcal{V}) \rightarrow \text{Clo}(\mathcal{W})$ . Say that  $\mathcal{V}$  is  **$\mathcal{L}$ -interpretable into**  $\mathcal{W}$ .

Define an equivalence relation by  $\mathcal{V} \equiv_{\mathcal{L}} \mathcal{W}$  if  $\mathcal{V} \leq_{\mathcal{L}} \mathcal{W} \leq_{\mathcal{L}} \mathcal{V}$ .

## Example

$$\mathcal{V}(\mathfrak{2}) \equiv_{\mathcal{L}} \mathcal{V}(\hat{\mathfrak{2}})$$

Let  $\mathcal{V}$  be a variety with signature  $\sigma$ .

Replace  $\sigma$  by  $\text{Clo}(\mathcal{V})$ , call the resulting variety  $\mathcal{V}_c$  so that  $\mathcal{V}_c \equiv \mathcal{V}$ .

Let  $\Sigma \subseteq \text{Th}(\mathcal{V}_c)$  be the subset of all identities not involving composition.

Let  $\mathcal{V}_{\mathcal{L}} = \text{Mod}(\Sigma)$ .

### Example

If  $\mathcal{V}$  has an identity  $t(s(x, y), x, y) \approx s(x, y)$ , then  $\mathcal{V}_c$  will have a new operation symbol  $f(x, y, z, w)$  and  $\text{Th}(\mathcal{V}_c)$  will have identities

- $t(s(x, y), x, y) \approx s(x, y)$ ,
- $f(x, y, z, w) \approx t(s(x, y), z, w)$ ,
- $f(x, y, x, y) \approx s(x, y)$ .

Out of these,  $\Sigma$  will contain **only** the last one.

### Example

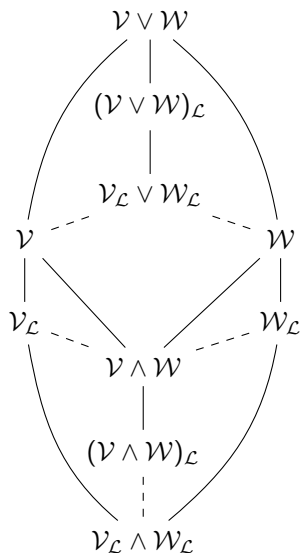
$(x \dashv x) \dashv y \approx (x \dashv x)$  in  $\hat{\mathcal{D}}$  becomes  $s(x, x, y) \approx x \dashv x$  in  $\mathcal{V}(\hat{\mathcal{D}})_{\mathcal{L}}$ .

## Theorem

Let  $\mathcal{V}$  and  $\mathcal{W}$  be varieties. Then

- $\mathcal{V} \equiv_{\mathcal{L}} \mathcal{V}_{\mathcal{L}}$ ,
- $\mathcal{V}_{\mathcal{L}} \leq \mathcal{W}_{\mathcal{L}}$  iff  $\mathcal{V} \leq_{\mathcal{L}} \mathcal{W}$ ,
- if  $\mathcal{V} \leq \mathcal{W}$  then  $\mathcal{V}_{\mathcal{L}} \leq \mathcal{W}_{\mathcal{L}}$ .
- $\mathcal{V}$  is defined by linear identities iff  $\mathcal{V} \equiv \mathcal{V}_{\mathcal{L}}$ .

Usual interpretability lattice:

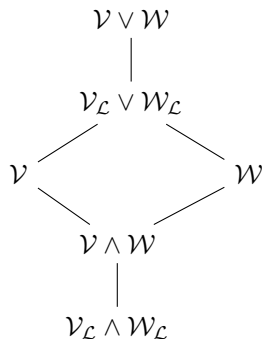


## Theorem

Let  $\mathcal{V}$  and  $\mathcal{W}$  be varieties. Then

- $\mathcal{V} \equiv_{\mathcal{L}} \mathcal{V}_{\mathcal{L}}$ ,
- $\mathcal{V}_{\mathcal{L}} \leq \mathcal{W}_{\mathcal{L}}$  iff  $\mathcal{V} \leq_{\mathcal{L}} \mathcal{W}$ ,
- if  $\mathcal{V} \leq \mathcal{W}$  then  $\mathcal{V}_{\mathcal{L}} \leq \mathcal{W}_{\mathcal{L}}$ .
- $\mathcal{V}$  is defined by linear identities iff  $\mathcal{V} \equiv \mathcal{V}_{\mathcal{L}}$ .

$\mathcal{L}$ -interpretability lattice:



## Corollary

Let  $\mathcal{V}$  be an idempotent variety. The following are equivalent.

- $\mathcal{V}$  does not have a cube term,
- $\mathcal{V} \leq_{\mathcal{L}} \mathcal{V}(\hat{\mathfrak{A}})$  in the  $\mathcal{L}$ -interpretability lattice.
- $\mathcal{V}_{\mathcal{L}} \leq \mathcal{V}(\hat{\mathfrak{A}})_{\mathcal{L}}$  in the interpretability lattice.

## Theorem (Sequeira; Barto, Oprsal, Pinsker)

If  $\mathcal{V}$  and  $\mathcal{W}$  are defined by linear identities and are  $\mathbb{A}$ -colorable, then  $\mathcal{V} \vee \mathcal{W}$  is  $\mathbb{A}$ -colorable as well.

## Corollary

If idempotent  $\mathcal{V}$  and  $\mathcal{W}$  are defined by linear identities and  $\mathcal{V} \vee \mathcal{W}$  has a cube term, then one of  $\mathcal{V}$  or  $\mathcal{W}$  does.

# Coloring varieties containing a cube term

- 1 Coloring
- 2 Clone  $\mathcal{L}$ -homomorphisms
- 3 Linear interpretability
- 4 Further directions

## Question

*Is having a cube term join-prime in the idempotent interpretability lattice?*

Answer is tentatively “yes”.

## Question

*Is there a characterization of when  $(\mathcal{V} \vee \mathcal{W})_{\mathcal{L}} \equiv \mathcal{V}_{\mathcal{L}} \vee \mathcal{W}_{\mathcal{L}}$ ?*

In the class of idempotent varieties satisfying  $(\mathcal{V} \vee \mathcal{W})_{\mathcal{L}} \equiv \mathcal{V}_{\mathcal{L}} \vee \mathcal{W}_{\mathcal{L}}$ , both being congruence modular and having a cube term are join prime.

## Question

*Does every  $\vee$ -prime filter in the  $\mathcal{L}$ -interpretability lattice come from a coloring characterization?*

## Question

*Is there a coloring characterization for  $k$ -permutability for fixed  $k$ ?*

We know that  $\mathcal{V}$  is not  $k$ -permutable for any  $k$  iff  $\mathcal{V}$  is  $(\{0, 1\}; \leq)$ -colorable.

Let  $\mathbb{P}_k = (\{1, \dots, k\}; \rightarrow)$  where  $\alpha \rightarrow \beta$  iff  $\alpha \leq \beta + 1$ .

Then  $\mathcal{V}$  is congruence  $(k - 1)$ -permutable iff it is not  $\mathbb{P}_k$ -colorable.



Let  $\hat{\mathbf{2}} = (\{0, 1\}; (\{0, 1\}^n \setminus \{1\}^n)_{n \in \omega})$  and  $\hat{\mathbf{2}} = \langle \{0, 1\}; \nrightarrow \rangle$ .

## Theorem

*For idempotent  $\mathcal{V}$  the following are equivalent.*

- $\mathcal{V}$  has no cube term,
- $\mathcal{V}$  is  $\hat{\mathbf{2}}$ -colorable,
- $\mathcal{V} \leq_{\mathcal{L}} \mathcal{V}(\hat{\mathbf{2}})$ ,
- $\mathcal{V}_{\mathcal{L}} \leq \mathcal{V}(\hat{\mathbf{2}})_{\mathcal{L}}$ .

*If  $\mathcal{V} = \mathcal{V}(\mathbb{A})$  for finite  $\mathbb{A}$ , then these are equivalent to*

- *there exists  $\mathbb{C} < \mathbb{D} \leq \mathbb{A}$  such that  $D^n \setminus (D \setminus C)^n \leq \mathbb{A}^n$  for all  $n$  (Markovic, Maroti, McKenzie).*

## Question

*Without repeating the proof of MMM, can we obtain the finite algebra result from the other results?*

# Conclusion

Let  $\hat{\mathbf{2}} = (\{0, 1\}; (\{0, 1\}^n \setminus \{1\}^n)_{n \in \omega})$  and  $\hat{\mathcal{2}} = \langle \{0, 1\}; \nrightarrow \rangle$ .

## Theorem

*For idempotent  $\mathcal{V}$  the following are equivalent.*

- $\mathcal{V}$  has no cube term,
- $\mathcal{V}$  is  $\hat{\mathbf{2}}$ -colorable,
- $\mathcal{V} \leq_{\mathcal{L}} \mathcal{V}(\hat{\mathcal{2}})$ ,
- $\mathcal{V}_{\mathcal{L}} \leq \mathcal{V}(\hat{\mathcal{2}})_{\mathcal{L}}$ .

Thank you.