

Supernilpotence Need Not Imply Nilpotence

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Let \mathbb{A} be an algebra and $\mathbb{R} \leq \mathbb{A}^4$.

$$\cancel{\begin{pmatrix} a \\ b \\ c \\ d \end{pmatrix} \in R}$$

instead, we will write

$$\begin{array}{ccc} a & \text{---} & b \\ | & & | \\ c & \text{---} & d \end{array} \in R$$

Let α_1, α_2 be congruences of \mathbb{A} . The (α_1, α_2) -**squares** of \mathbb{A} are

$$M(\alpha_1, \alpha_2) := \text{Sg}_{\mathbb{A}^4} \left\{ \begin{array}{ccc} a & \text{---} & b \\ | & & | \\ a & \text{---} & b \end{array}, \begin{array}{ccc} c & \text{---} & c \\ | & & | \\ d & \text{---} & d \end{array} \mid (a, b) \in \alpha_1, (c, d) \in \alpha_2 \right\}$$

Observe

$$\begin{array}{ccc} a & \text{---} & b \\ | & & | \\ c & \text{---} & d \end{array} \in M(\alpha_1, \alpha_2) \quad \Longrightarrow \quad \alpha_2 \left(\begin{array}{ccc} & \overset{\alpha_1}{\text{---}} & \\ a & \text{---} & b \\ | & & | \\ c & \text{---} & d \\ & \underset{\alpha_1}{\text{---}} & \end{array} \right) \alpha_2$$

$$M(\alpha_1, \alpha_2) := \text{Sg}_{\mathbb{A}^4} \left\{ \begin{array}{cc|cc} a & \text{---} & b & & c & \text{---} & c \\ | & & | & & | & & | \\ a & \text{---} & b & & d & \text{---} & d \end{array} \mid (a, b) \in \alpha_1, (c, d) \in \alpha_2 \right\}$$

α_1 **centralizes** α_2 modulo δ if ...

[write $C(\alpha_1, \alpha_2; \delta)$]

$$\forall \begin{array}{cc|cc} a & \text{---} & b & & c & \text{---} & d \\ | & & | & & | & & | \\ c & \text{---} & d & & & & \end{array} \in M(\alpha_1, \alpha_2), \quad \delta \left(\begin{array}{cc|cc} a & \text{---} & b & & c & \text{---} & d \\ | & & | & & | & & | \\ c & \text{---} & d & & & & \end{array} \right) \Rightarrow \begin{array}{cc|cc} a & \text{---} & b & & c & \text{---} & d \\ | & & | & & | & & | \\ c & \text{---} & d & & & & \end{array} \delta$$

The **commutator** of α_1 and α_2 is the smallest δ satisfying the above,

$$[\alpha_1, \alpha_2] := \bigwedge_{C(\alpha_1, \alpha_2; \delta)} \delta$$

Can we generalize this to a higher arity commutator?

Let \mathbb{A} be an algebra and $\mathbb{R} \leq \mathbb{A}^{2^3}$.

$$\left(\begin{array}{c} a \\ \vdots \\ h \end{array} \right) \in R$$

instead, we will write

$$\begin{array}{ccccc} a & \text{---} & b & & \\ & \diagdown & e & \text{---} & \diagup & f \\ & & & \text{---} & & \\ c & & & \text{---} & & d \\ & \diagup & & & \diagdown & \\ & & g & \text{---} & & h \end{array} \in R$$

Let $\alpha_1, \alpha_2, \alpha_3$ be congruences of \mathbb{A} . The $(\alpha_1, \alpha_2, \alpha_3)$ -**cubes** of \mathbb{A} are

$$M(\alpha_1, \alpha_2, \alpha_3) := \text{Sg}_{\mathbb{A}^{2^3}} \left\{ \begin{array}{c} \begin{array}{ccccc} a & \text{---} & b & & \\ & \diagdown & a & \text{---} & \diagup & b \\ & & & \text{---} & & \\ a & & & \text{---} & & b \\ & \diagup & & & \diagdown & \\ & & a & \text{---} & & b \end{array}, \begin{array}{ccccc} c & \text{---} & c & & \\ & \diagdown & c & \text{---} & \diagup & c \\ & & & \text{---} & & \\ d & & & \text{---} & & d \\ & \diagup & & & \diagdown & \\ & & d & \text{---} & & d \end{array}, \begin{array}{ccccc} e & \text{---} & e & & \\ & \diagdown & f & \text{---} & \diagup & f \\ & & & \text{---} & & \\ e & & & \text{---} & & e \\ & \diagup & & & \diagdown & \\ & & f & \text{---} & & f \end{array} \mid \begin{array}{l} (a, b) \in \alpha_1, \\ (c, d) \in \alpha_2, \\ (e, f) \in \alpha_3 \end{array} \right\}$$

Observe

$$\begin{array}{ccccc} a & \text{---} & b & & \\ & \diagdown & e & \text{---} & \diagup & f \\ & & & \text{---} & & \\ c & & & \text{---} & & d \\ & \diagup & & & \diagdown & \\ & & g & \text{---} & & h \end{array} \in M(\alpha_1, \alpha_2, \alpha_3) \implies \begin{array}{c} \alpha_2 \\ \alpha_1 \\ \alpha_3 \end{array} \text{ in } \begin{array}{ccccc} a & \text{---} & b & & \\ & \diagdown & e & \text{---} & \diagup & f \\ & & & \text{---} & & \\ c & & & \text{---} & & d \\ & \diagup & & & \diagdown & \\ & & g & \text{---} & & h \end{array}$$

$$M(\alpha_1, \alpha_2, \alpha_3) := \text{Sg}_{\mathbb{A}^{2^3}} \left\{ \begin{array}{ccc} \begin{array}{c} a \text{ --- } b \\ | \quad \backslash \quad \backslash \\ a \text{ --- } a \text{ --- } b \\ | \quad \backslash \quad \backslash \\ a \text{ --- } b \end{array}, & \begin{array}{c} c \text{ --- } c \\ | \quad \backslash \quad \backslash \\ d \text{ --- } d \\ | \quad \backslash \quad \backslash \\ d \text{ --- } d \end{array}, & \begin{array}{c} e \text{ --- } e \\ | \quad \backslash \quad \backslash \\ f \text{ --- } f \\ | \quad \backslash \quad \backslash \\ f \text{ --- } f \end{array} \quad \left. \begin{array}{l} (a, b) \in \alpha_1, \\ (c, d) \in \alpha_2, \\ (e, f) \in \alpha_3 \end{array} \right\}$$

α_1, α_2 **centralize** α_3 modulo δ if ...

[write $C(\alpha_1, \alpha_2, \alpha_3; \delta)$]

$$\forall \begin{array}{c} a \text{ --- } b \\ | \quad \backslash \quad \backslash \\ c \text{ --- } e \text{ --- } f \\ | \quad \backslash \quad \backslash \\ g \text{ --- } d \text{ --- } h \end{array} \in M(\alpha_1, \alpha_2, \alpha_3), \quad \begin{array}{c} a \text{ --- } b \\ | \quad \backslash \quad \backslash \\ \delta \text{ --- } e \text{ --- } f \\ | \quad \backslash \quad \backslash \\ \delta \text{ --- } c \text{ --- } d \\ | \quad \backslash \quad \backslash \\ \delta \text{ --- } g \text{ --- } h \end{array} \implies \begin{array}{c} a \text{ --- } b \text{ --- } \delta \\ | \quad \backslash \quad \backslash \\ c \text{ --- } e \text{ --- } f \\ | \quad \backslash \quad \backslash \\ g \text{ --- } d \text{ --- } h \end{array}$$

The **3-dim commutator** of $\alpha_1, \alpha_2, \alpha_3$ is the smallest δ satisfying the above,

$$[\alpha_1, \alpha_2, \alpha_3] := \bigwedge_{C(\alpha_1, \alpha_2, \alpha_3; \delta)} \delta$$

Can we generalize this to a higher arity commutator?

Yes, even for arbitrary $n!$

A general definition of the commutator yields general notions of abelianness, solvability, and nilpotence.

- \mathbb{A} is **abelian** if $[\mathbb{1}, \mathbb{1}] = \mathbb{0}$.
- Define the **derived series**,

$$[\alpha]_0 := \alpha \qquad [\alpha]_{n+1} := [[\alpha]_n, [\alpha]_n].$$

\mathbb{A} is **m -solvable** if $[\mathbb{1}]_m = \mathbb{0}$ for some m .

- define the **lower central series**,

$$(\alpha)_0 := \alpha \qquad (\alpha)_{n+1} := [\alpha, (\alpha)_n].$$

\mathbb{A} is (left) **m -nilpotent** if $(\mathbb{1})_m = \mathbb{0}$ for some m .

- The higher dimensional commutator has its own series,

$$\alpha \geq [\alpha, \alpha] \geq [\alpha, \alpha, \alpha] \geq \cdots \geq [\alpha, \dots, \alpha] \geq \dots$$

\mathbb{A} is **$(m+1)$ -supernilpotent** if $[\mathbb{1}, \dots, \mathbb{1}] = \mathbb{0}$ (m -ary)

($\mathbb{1}$ is the universal congruence and $\mathbb{0}$ is the trivial congruence)

For an algebra \mathbb{A} we have general notions of

- abelianness
- solvability
- nilpotence
- supernilpotence

How are these related?

How are nilpotence and supernilpotence related?

Does supernilpotence imply nilpotence?

Theorem (Moorhead)

If \mathbb{A} satisfies a non-trivial idempotent equational condition, then “Yes”.

Theorem (Kearnes, Szendrei)

If \mathbb{A} is finite, then “Yes”.

Theorem (M+M)

Supernilpotence does not imply nilpotence in general.

Idea: Construct \mathbb{A} as simply as possible.

- Make \mathbb{A} 2-supernilpotent: $[\mathbb{1}, \mathbb{1}, \mathbb{1}] = \mathbb{0}$.
- When $[\cdot, \cdot]$ is not symmetric, there are many different notions of nilpotence. All of them imply solvability.
- Make \mathbb{A} not solvable.

What does $[\mathbb{1}, \mathbb{1}, \mathbb{1}] = \mathbb{0}$ mean?

$$M(\mathbb{1}, \mathbb{1}, \mathbb{1}) := \text{Sg}_{\mathbb{A}^{2^3}} \left\{ \begin{array}{c} a \text{ --- } b \\ | \quad \diagdown \quad | \quad \diagdown \\ a \quad \text{---} \quad a \text{ ---} \quad b \\ | \quad \diagup \quad | \quad \diagup \\ a \quad \text{---} \quad b \end{array}, \begin{array}{c} a \text{ --- } a \\ | \quad \diagdown \quad | \quad \diagdown \\ a \quad \text{---} \quad a \\ | \quad \diagup \quad | \quad \diagup \\ b \quad \text{---} \quad b \end{array}, \begin{array}{c} a \text{ --- } a \\ | \quad \diagdown \quad | \quad \diagdown \\ a \quad \text{---} \quad b \\ | \quad \diagup \quad | \quad \diagup \\ a \quad \text{---} \quad a \\ | \quad \diagdown \quad | \quad \diagdown \\ b \quad \text{---} \quad b \end{array} \mid a, b \in A \right\}$$

$$\begin{array}{c} a \text{ --- } b \\ | \quad \diagdown \quad | \quad \diagdown \\ a \quad \text{---} \quad f \\ | \quad \diagup \quad | \quad \diagup \\ c \quad \text{---} \quad d \\ | \quad \diagdown \quad | \quad \diagdown \\ c \quad \text{---} \quad d \end{array} \in M(\mathbb{1}, \mathbb{1}, \mathbb{1}) \quad \implies \quad b = f$$

We will need to carefully analyze the generation of $M(\mathbb{1}, \mathbb{1}, \mathbb{1})$.

$$M(\mathbb{1}, \mathbb{1}, \mathbb{1}) := \text{Sg}_{\mathbb{A}^{23}} \left\{ \begin{array}{c} a \text{ --- } b \\ | \quad \diagdown \\ | \quad a \text{ --- } b \\ | \quad | \quad \diagdown \\ a \text{ --- } b \\ | \quad | \quad \diagdown \\ a \text{ --- } b \end{array} , \begin{array}{c} a \text{ --- } a \\ | \quad \diagdown \\ | \quad a \text{ --- } a \\ | \quad | \quad \diagdown \\ b \text{ --- } b \\ | \quad | \quad \diagdown \\ b \text{ --- } b \end{array} , \begin{array}{c} a \text{ --- } a \\ | \quad \diagdown \\ | \quad b \text{ --- } b \\ | \quad | \quad \diagdown \\ a \text{ --- } a \\ | \quad | \quad \diagdown \\ b \text{ --- } b \end{array} \mid a, b \in A \right\}$$

$$[\mathbb{1}, \mathbb{1}, \mathbb{1}] = 0 \quad \text{iff} \quad \left[\begin{array}{c} a \text{ --- } b \\ | \quad \diagdown \\ | \quad a \text{ --- } f \\ | \quad | \quad \diagdown \\ c \text{ --- } d \\ | \quad | \quad \diagdown \\ c \text{ --- } d \end{array} \in M(\mathbb{1}, \mathbb{1}, \mathbb{1}) \implies b = f \right] \quad (\dagger)$$

What is the simplest way to ensure $[\mathbb{1}, \mathbb{1}, \mathbb{1}] = 0$?

- consider $\mathbb{A} = \langle A ; t(x, y) \rangle$ where $t(x, y)$ is injective (infinite \mathbb{A})
- generating $M(\mathbb{1}, \mathbb{1}, \mathbb{1})$ in layers, look for first failure of (\dagger)

$$t \left(\begin{array}{c} a_1 \text{ --- } b_1 \\ | \quad \diagdown \\ | \quad a_1 \text{ --- } f_1 \\ | \quad | \quad \diagdown \\ c_1 \text{ --- } d_1 \\ | \quad | \quad \diagdown \\ c_1 \text{ --- } d_1 \end{array} , \begin{array}{c} a_2 \text{ --- } b_2 \\ | \quad \diagdown \\ | \quad a_2 \text{ --- } f_2 \\ | \quad | \quad \diagdown \\ c_2 \text{ --- } d_2 \\ | \quad | \quad \diagdown \\ c_2 \text{ --- } d_2 \end{array} \right) = \begin{array}{c} a \text{ --- } b \\ | \quad \diagdown \\ | \quad a \text{ --- } f \\ | \quad | \quad \diagdown \\ c \text{ --- } d \\ | \quad | \quad \diagdown \\ c \text{ --- } d \end{array} \quad (\text{all in } M(\mathbb{1}, \mathbb{1}, \mathbb{1}))$$

- \mathbb{A} is 2-supernilpotent! \mathbb{A} is also nilpotent!

$$M(\mathbb{1}, \mathbb{1}, \mathbb{1}) := \text{Sg}_{\mathbb{A}^{23}} \left\{ \begin{array}{c} a \text{ --- } b \\ | \quad \diagdown \\ | \quad a \text{ --- } b \\ | \quad | \quad \diagdown \\ a \text{ --- } b \\ | \quad | \quad \diagdown \\ a \text{ --- } b \end{array} , \begin{array}{c} a \text{ --- } a \\ | \quad \diagdown \\ | \quad a \text{ --- } a \\ | \quad | \quad \diagdown \\ b \text{ --- } b \\ | \quad | \quad \diagdown \\ b \text{ --- } b \end{array} , \begin{array}{c} a \text{ --- } a \\ | \quad \diagdown \\ | \quad b \text{ --- } b \\ | \quad | \quad \diagdown \\ a \text{ --- } a \\ | \quad | \quad \diagdown \\ b \text{ --- } b \end{array} \mid a, b \in A \right\}$$

$$[\mathbb{1}, \mathbb{1}, \mathbb{1}] = 0 \quad \text{iff} \quad \left[\begin{array}{c} a \text{ --- } b \\ | \quad \diagdown \\ | \quad a \text{ --- } f \\ | \quad | \quad \diagdown \\ c \text{ --- } d \\ | \quad | \quad \diagdown \\ c \text{ --- } d \end{array} \in M(\mathbb{1}, \mathbb{1}, \mathbb{1}) \implies b = f \right] \quad (\dagger)$$

What is the simplest way to ensure $[\mathbb{1}, \mathbb{1}, \mathbb{1}] = 0$?

- consider $\mathbb{A} = \langle A ; t(x, y) \rangle$ where $t(x, y)$ is injective (infinite \mathbb{A})
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$$t \left(\begin{array}{c} a_1 \text{ --- } b_1 \\ | \quad \diagdown \\ | \quad a_1 \text{ --- } b_1 \\ | \quad | \quad \diagdown \\ c_1 \text{ --- } d_1 \\ | \quad | \quad \diagdown \\ c_1 \text{ --- } d_1 \end{array} , \begin{array}{c} a_2 \text{ --- } b_2 \\ | \quad \diagdown \\ | \quad a_2 \text{ --- } b_2 \\ | \quad | \quad \diagdown \\ c_2 \text{ --- } d_2 \\ | \quad | \quad \diagdown \\ c_2 \text{ --- } d_2 \end{array} \right) = \begin{array}{c} a \text{ --- } b \\ | \quad \diagdown \\ | \quad a \text{ --- } b \\ | \quad | \quad \diagdown \\ c \text{ --- } d \\ | \quad | \quad \diagdown \\ c \text{ --- } d \end{array} \quad (\text{all in } M(\mathbb{1}, \mathbb{1}, \mathbb{1}))$$

- \mathbb{A} is 2-supernilpotent! \mathbb{A} is also nilpotent!

Idea:

- $\mathbb{A} = \langle A ; t(x, y) \rangle$, injective $t(x, y)$
- carefully redefine $t(x, y)$ to sabotage injectivity and nilpotence
- be clever so that $t(x, y)$ is injective enough to ensure $[1, 1, 1] = 0$

Define

$$O := \{\sigma_i^j \mid i, j \in \mathbb{N}\}, \quad R := \{r_i^j \mid i, j \in \mathbb{N}\}, \quad A := O \cup R \cup \mathbb{N}.$$

Fix an injection $s : A^2 \rightarrow \mathbb{N}$. Define

$$t \left(\begin{array}{cc} r_{4i}^j - r_{4i+2}^j & r_{4i}^j - r_{4i}^j \\ | & | \\ r_{4i}^j - r_{4i+2}^j & r_{4i+2}^j - r_{4i+2}^j \end{array} , \begin{array}{cc} | & | \\ | & | \end{array} \right) := \begin{array}{c} \sigma_i^j - r_i^{j+1} \\ | \\ \sigma_i^j - r_{i+1}^{j+1} \end{array}$$

Otherwise, define $t(x, y) := s(x, y)$. Let $\mathbb{A} = \langle A ; t(x, y) \rangle$.

Observe: $t(x, y)$ is injective except when the output is in O .

Define $A = O \cup R \cup \mathbb{N}$ and

$$t \left(\begin{array}{cc} r_{4i}^j - r_{4i+2}^j & r_{4i}^j - r_{4i}^j \\ | & | \\ r_{4i}^j - r_{4i+2}^j & r_{4i+2}^j - r_{4i+2}^j \end{array} \right) := \begin{array}{cc} \sigma_i^j - r_i^{j+1} \\ | & | \\ \sigma_i^j - r_{i+1}^{j+1} \end{array},$$

otherwise define $t(x, y) := s(x, y)$ for some injection $s : A^2 \rightarrow \mathbb{N}$.

\mathbb{A} is not solvable

- Derived series: $[\alpha]_0 := \alpha$, $[\alpha]_{n+1} := [[\alpha]_n, [\alpha]_n]$.
- $R^j := \{r_i^j \mid i, j \in \mathbb{N}\}$ is contained in a $[\mathbb{1}]_j$ -class (induction on j).
- Thus $[\mathbb{1}]_j \neq \emptyset$ for any j , so \mathbb{A} is not solvable.

Define $A = O \cup R \cup \mathbb{N}$ and

$$t \left(\begin{array}{cc|cc} r_{4i}^j - r_{4i+2}^j & r_{4i}^j - r_{4i}^j & & \\ | & | & & \\ r_{4i}^j - r_{4i+2}^j & r_{4i+2}^j - r_{4i+2}^j & & \end{array} , \begin{array}{cc|cc} r_{4i}^j - r_{4i}^j & & & \\ | & & & \\ r_{4i+2}^j - r_{4i+2}^j & & & \end{array} \right) := \begin{array}{cc|cc} \sigma_i^j - r_i^{j+1} & & & \\ | & & & \\ \sigma_i^j - r_{i+1}^{j+1} & & & \end{array} ,$$

otherwise define $t(x, y) := s(x, y)$ for some injection $s : A^2 \rightarrow \mathbb{N}$.

Generating $M(\mathbb{1}, \mathbb{1}, \mathbb{1})$ in layers, look for first instance of

$$t \left(\begin{array}{cc|cc} a_1 - b_1 & & a_2 - b_2 & \\ | & e_1 - f_1 & | & e_2 - f_2 \\ c_1 - d_1 & & c_2 - d_2 & \\ | & g_1 - h_1 & | & g_2 - h_2 \end{array} , \begin{array}{cc|cc} a - b & & & \\ | & a - f & & \\ c - d & & & \\ | & c - d & & \end{array} \right) = \begin{array}{cc|cc} a - b & & & \\ | & a - f & & \\ c - d & & & \\ | & c - d & & \end{array} , \quad b \neq f$$

Observations

- If $t(a, b) = t(c, d)$ then $a = c$.
- Argument squares: if 3 back-to-front edges are equal, then the 4th is too.
- $b_2 \neq f_2$, otherwise $b = f$. At most 2 back-to-front edges in the second argument cube are equal.

The rest of the argument is technical ...

Define $A = O \cup R \cup \mathbb{N}$ and

$$t \left(\begin{array}{cc} r_{4i}^j - r_{4i+2}^j & r_{4i}^j - r_{4i}^j \\ | & | \\ r_{4i}^j - r_{4i+2}^j & r_{4i+2}^j - r_{4i+2}^j \end{array}, \begin{array}{cc} r_{4i}^j - r_{4i}^j & \\ | & \\ r_{4i+2}^j - r_{4i+2}^j & \end{array} \right) := \begin{array}{cc} \sigma_i^j - r_i^{j+1} & \\ | & | \\ \sigma_i^j - r_{i+1}^{j+1} & \end{array},$$

otherwise define $t(x, y) := s(x, y)$ for some injection $s : A^2 \rightarrow \mathbb{N}$.

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$$t \left(\begin{array}{cc} a_1 - b_1 & a_2 - b_2 \\ | & | \\ c_1 - d_1 & c_2 - d_2 \\ | & | \\ c_1 - d_1 & g_2 - h_2 \end{array}, \begin{array}{cc} a_1 - b_1 & \\ | & \\ c_1 - d_1 & \\ | & \\ c_1 - d_1 & \end{array} \right) = \begin{array}{cc} a - b & \\ | & | \\ c - d & \\ | & | \\ c - d & \end{array}, \quad b \neq f$$

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Theorem (M+M)

The algebra \mathbb{A} is 3-supernilpotent but not solvable.

- nilpotence and solvability can be generalized using the n -ary commutator
- call the generalized notions n -dimensional nilpotence and solvability
- a similar construction yields $\mathbb{A}_n := \langle A ; t(x_1, \dots, x_n) \rangle$.

Theorem (M+M)

The algebra \mathbb{A}_n is $(n + 1)$ -supernilpotent but not n -dimensional solvable.

Supernilpotence Need Not Imply Nilpotence

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Thank you for your attention.